

# An automaton describing left burnable configurations in the sandpile model on the ladder graph

Henri Derycke

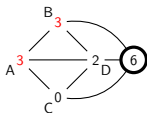
LaBRI

ICGT 2018, July 9-13, Lyon



université  
de BORDEAUX

# The sandpile model

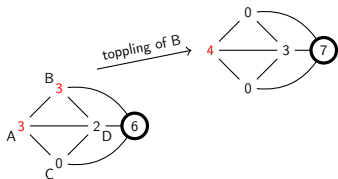


## The sandpile model

$G = (V, E)$  simple connected graph with a special vertex  $S \in V$  called the *sink*

- ▶ Configuration : assignment  $u : V \rightarrow \mathbf{N}$  of grains on the vertices
- ▶ A vertex distinct from the sink is *unstable* if  $u(x) \geq \deg(x)$
- ▶ Toppling : a vertex sends one grain to each of its neighbors
- ▶ Stabilization : While there is an unstable vertex, we topple it. Given  $u$ , we note  $\text{stab}(u)$  the result.

# The sandpile model

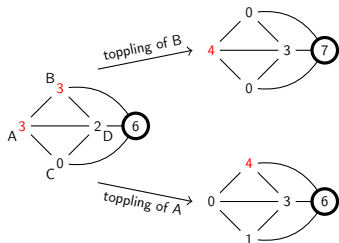


## The sandpile model

$G = (V, E)$  simple connected graph with a special vertex  $S \in V$  called the *sink*

- ▶ Configuration : assignment  $u : V \rightarrow \mathbf{N}$  of grains on the vertices
- ▶ A vertex distinct from the sink is *unstable* if  $u(x) \geq \deg(x)$
- ▶ Toppling : a vertex sends one grain to each of its neighbors
- ▶ Stabilization : While there is an unstable vertex, we topple it. Given  $u$ , we note  $\text{stab}(u)$  the result.

# The sandpile model

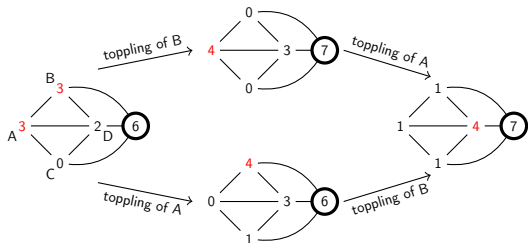


## The sandpile model

$G = (V, E)$  simple connected graph with a special vertex  $S \in V$  called the *sink*

- ▶ Configuration : assignment  $u : V \rightarrow \mathbf{N}$  of grains on the vertices
- ▶ A vertex distinct from the sink is *unstable* if  $u(x) \geq \deg(x)$
- ▶ Toppling : a vertex sends one grain to each of its neighbors
- ▶ Stabilization : While there is an unstable vertex, we topple it. Given  $u$ , we note  $\text{stab}(u)$  the result.

# The sandpile model

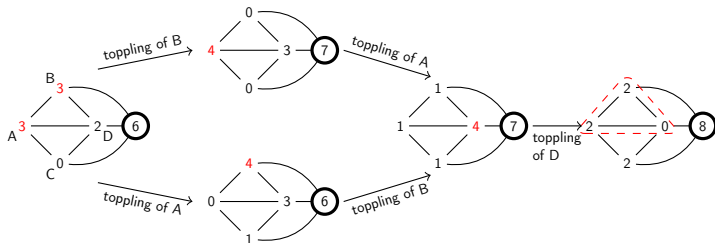


## The sandpile model

$G = (V, E)$  simple connected graph with a special vertex  $S \in V$  called the *sink*

- ▶ Configuration : assignment  $u : V \rightarrow \mathbf{N}$  of grains on the vertices
- ▶ A vertex distinct from the sink is *unstable* if  $u(x) \geq \deg(x)$
- ▶ Toppling : a vertex sends one grain to each of its neighbors
- ▶ Stabilization : While there is an unstable vertex, we topple it. Given  $u$ , we note  $\text{stab}(u)$  the result.

# The sandpile model



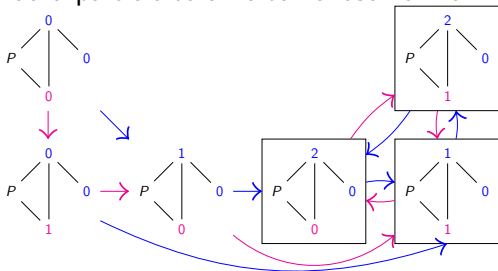
## The sandpile model

$G = (V, E)$  simple connected graph with a special vertex  $S \in V$  called the *sink*

- ▶ Configuration : assignment  $u : V \rightarrow \mathbf{N}$  of grains on the vertices
- ▶ A vertex distinct from the sink is *unstable* if  $u(x) \geq \deg(x)$
- ▶ Toppling : a vertex sends one grain to each of its neighbors
- ▶ Stabilization : While there is an unstable vertex, we topple it. Given  $u$ , we note  $\text{stab}(u)$  the result.

# Markov Chain for $G = (V \cup \{S\}, E)$

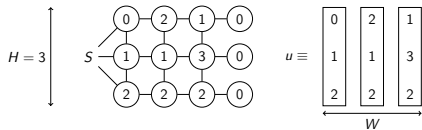
- ▶ States : stable configurations on  $G$
- ▶ Transition : Add a particle to a vertex chosen uniformly and stabilize



- ▶ The recurrent states are called recurrent configurations.
- ▶ The stationary distribution is uniform on the recurrent configurations.

**Dhar Criterion** A stable configuration is recurrent if and only if adding a grain to each neighbor of the sink, and stabilizing result to the same configuration. (fixed point)

# Configurations on left rooted square grid graphs

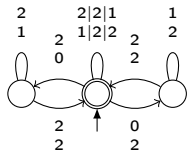


A configuration  $u$  may be read as a word of size  $W$  on the alphabet of columns.

Járai, Lyons 07

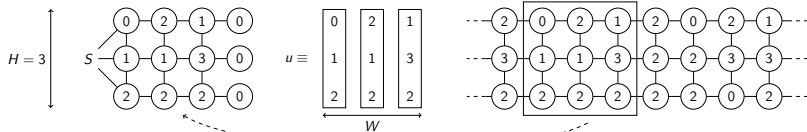
The recurrent configurations of the left-rooted square grid graphs are recognized by a finite automaton on the columns.

For height 2, they found the following automaton.





# Configurations on left rooted square grid graphs

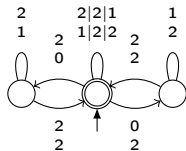


A configuration  $u$  may be read as a word of size  $W$  on the alphabet of columns.

Járai, Lyons 07

The recurrent configurations of the left-rooted square grid graphs are recognized by a finite automaton on the columns.

For height 2, they found the following automaton.



## Some results on these automata of height $H \geq 2$

From J arai and Lyons, the number of states is roughly bounded by  $(2^H)^{2^H}$ .

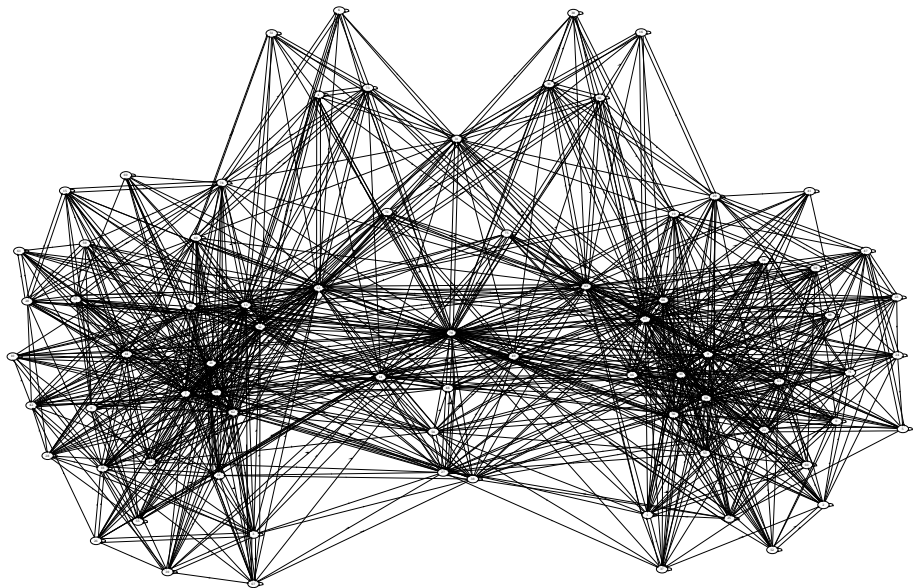
**Conjecture from Gamlin (Thesis 2015)**  $\exp cH$  states may be enough.

Similar to an automaton for spanning trees (eg. dominoes tiling) but optimistic on the non local transformation to recurrent configurations.

### My results

- ▶ Efficient construction of an automaton.
- ▶ From  $\alpha^H$  to  $\beta^{H \log H}$  states.
- ▶ The lower bound is from a bijection between a subset of the states and separable permutations (avoiding patterns 2413 and 3142).
- ▶ Explicit automata for heights 3 and 4





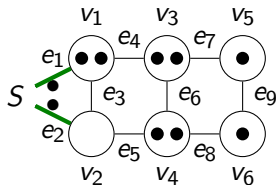
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3, e_4, v_3, e_6, v_4, e_5, v_2, e_7, v_5, e_8, v_6, e_9$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
     untreated incident edges as pending edges.

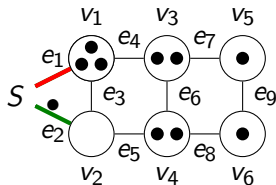
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1$ ,



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
     untreated incident edges as pending edges.

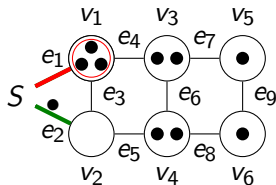
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1$ ,



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
     untreated incident edges as pending edges.

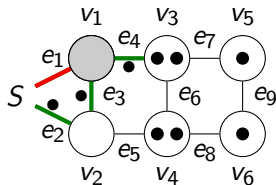
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1,$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
 Get the smallest pending edge  
 Process the grain(s) on the edge  
 If a vertex become unstable, topple it and mark its  
 untreated incident edges as pending edges.



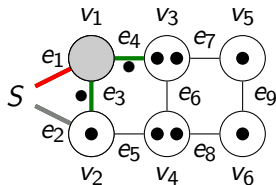
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2,$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
   untreated incident edges as pending edges.

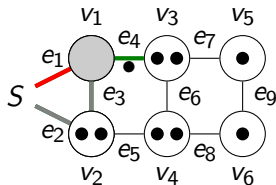
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3,$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
   untreated incident edges as pending edges.

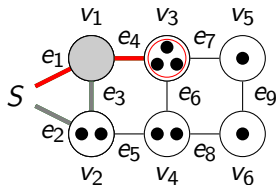
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3, e_4,$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
   untreated incident edges as pending edges.

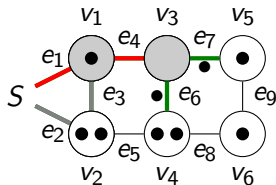
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3, e_4, v_3,$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
   untreated incident edges as pending edges.

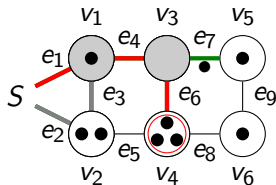
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3, e_4, v_3, e_6,$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
   untreated incident edges as pending edges.

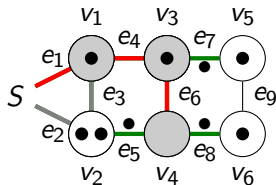
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3, e_4, v_3, e_6, v_4,$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
   untreated incident edges as pending edges.

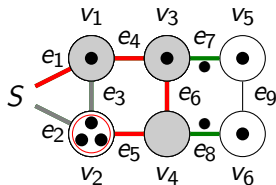
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3, e_4, v_3, e_6, v_4, e_5,$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
   untreated incident edges as pending edges.

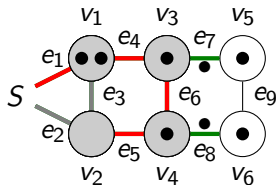
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3, e_4, v_3, e_6, v_4, e_5, v_2,$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
   untreated incident edges as pending edges.



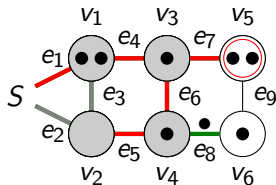
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3, e_4, v_3, e_6, v_4, e_5, v_2, e_7,$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
   untreated incident edges as pending edges.

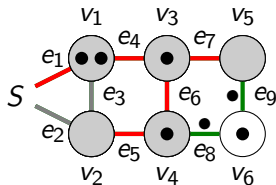
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3, e_4, v_3, e_6, v_4, e_5, v_2, e_7, v_5,$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
 Get the smallest pending edge  
 Process the grain(s) on the edge  
 If a vertex become unstable, topple it and mark its  
 untreated incident edges as pending edges.

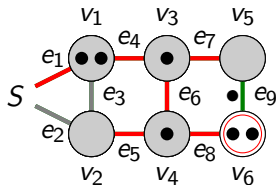
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3, e_4, v_3, e_6, v_4, e_5, v_2, e_7, v_5, e_8$ ,



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
     untreated incident edges as pending edges.

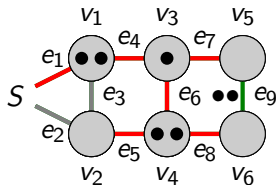
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3, e_4, v_3, e_6, v_4, e_5, v_2, e_7, v_5, e_8, v_6, e_9$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
   untreated incident edges as pending edges.

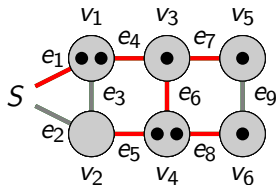
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3, e_4, v_3, e_6, v_4, e_5, v_2, e_7, v_5, e_8, v_6, e_9$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
   Get the smallest pending edge  
   Process the grain(s) on the edge  
   If a vertex become unstable, topple it and mark its  
   untreated incident edges as pending edges.

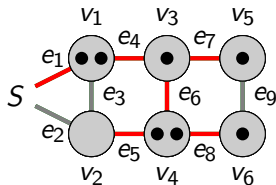
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3, e_4, v_3, e_6, v_4, e_5, v_2, e_7, v_5, e_8, v_6, e_9$



Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
 Get the smallest pending edge  
 Process the grain(s) on the edge  
 If a vertex become unstable, topple it and mark its  
 untreated incident edges as pending edges.

- ▶ The spanning tree is encoded by the red edges that precede vertices in the decreasing edge-vertex traversal.
- ▶ For any vertex, the number of incident edges that are treated after the vertex is equal to the number of grains in the recurrent configuration.

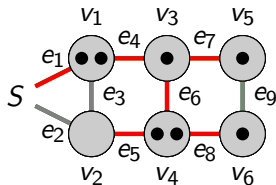
# Bijections between recurrent configurations and spanning trees

Some bijections : Dhar/Majumdar 92, Cori/Le Borgne 03, Bernardi 06...

We choose an order on the edges : from left to right and from top to bottom.

**Bijection CLB** : (from a recurrent configuration)

Decreasing traversals :  $e_1, v_1, e_2, e_3, e_4, v_3, e_6, v_4, e_5, v_2, e_7, e_8, v_6, e_9, v_5$

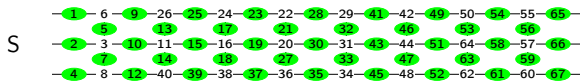


Mark edges incident to the sink as pending edges.  
 While there is a pending edge  
 Get the smallest pending edge  
 Process the grain(s) on the edge  
 If a vertex become unstable, topple it and mark its  
 untreated incident edges as pending edges.

- ▶ The spanning tree is encoded by the red edges that precede vertices in the decreasing edge-vertex traversal.
- ▶ For any vertex, the number of incident edges that are treated after the vertex is equal to the number of grains in the recurrent configuration.

Edge — X —

Vertex — y —

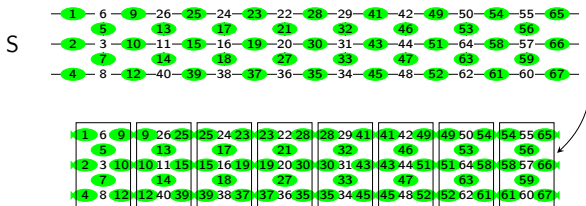




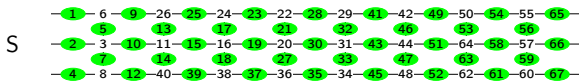
Edge —X—

Vertex —y—

- Decomposition

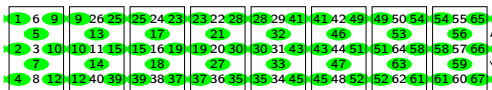


Edge -x-

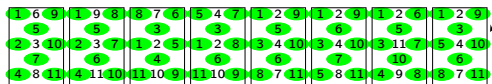


Vertex -y-

- Decomposition

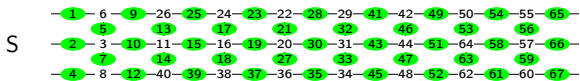


- Standardization  
( $\rightarrow$  finite alphabet)

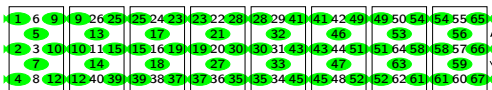


Edge -X-

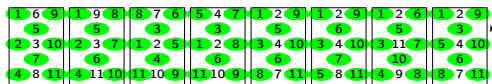
Vertex -y-



- Decomposition

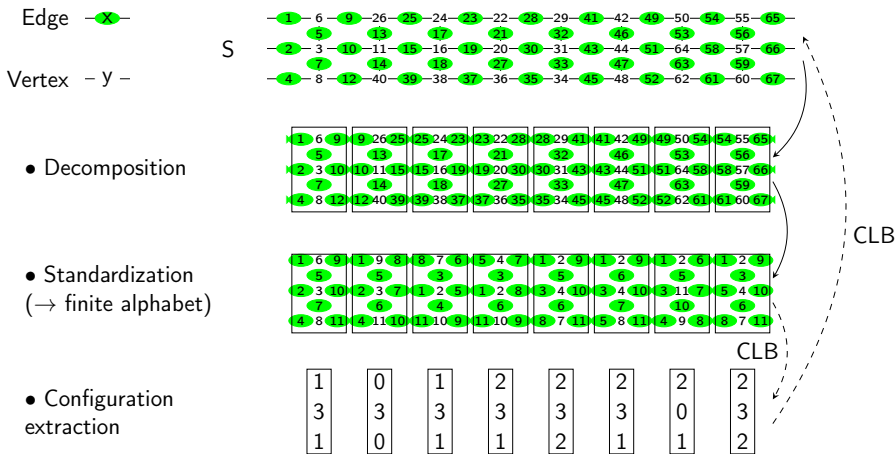


- Standardization  
( $\rightarrow$  finite alphabet)



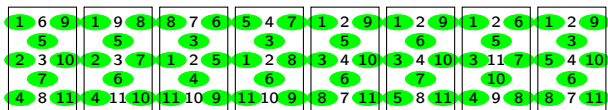
## Proposition

The edge-vertex decreasing traversals are in bijection with their standardized decompositions (in induced orders).



## Proposition

The edge-vertex decreasing traversals are in bijection with their standardized decompositions (in induced orders).



A decomposition is a word on the alphabet of induced orders.

- $e_1$   $v_1$   $e_6$  ▶ Initial induced order :  $e_1(v_1?)e_2(v_2?)e_3(v_3?)\dots$
- $e_4$  ▶ Sequence of consecutive pairwise compatible induced orders :  
 $(o, o')$  verify  $T(o, o')$
- $e_2$   $v_2$   $e_7$
- $e_5$
- $e_3$   $v_3$   $e_8$  ▶ Final induced order :  $\dots e_6 e_7 e_8$

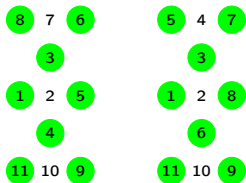
Automaton on induced orders

- ▶ States : induced orders
- ▶ Alphabet : induced orders
- ▶ Transition :  $o \xrightarrow{o'} o'$  iff  $o'$  is right-compatible to  $o$

**Prop**  $\#stats \leq (4H - 1)!$

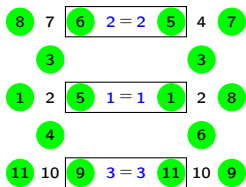
We have an automaton recognizing recurrent configurations

# Right compatibility relation



## Necessary conditions

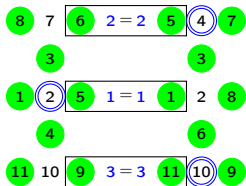
# Right compatibility relation



## Necessary conditions

- ▶ the orders on common edges are the same

# Right compatibility relation

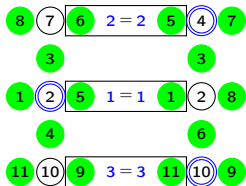


## Necessary conditions

- ▶ the orders on common edges are the same
- ▶ horizontal edges are after at least one of their endpoints



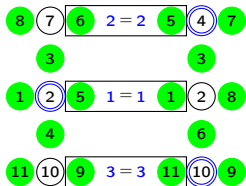
# Right compatibility relation



## Necessary conditions

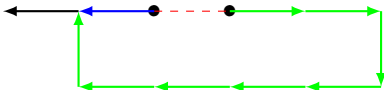
- ▶ the orders on common edges are the same
- ▶ horizontal edges are after at least one of their endpoints
- ▶ horizontal edges are before at least one of their endpoints : (the maximal edge of any cycle is vertical)

# Right compatibility relation

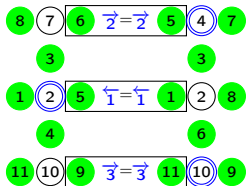


## Necessary conditions

- ▶ the orders on common edges are the same
- ▶ horizontal edges are after at least one of their endpoints
- ▶ horizontal edges are before at least one of their endpoints : (the maximal edge of any cycle is vertical)

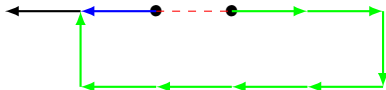


# Right compatibility relation

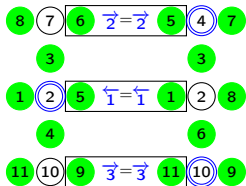


## Necessary conditions

- ▶ the orders on common edges are the same
- ▶ horizontal edges are after at least one of their endpoints
- ▶ horizontal edges are before at least one of their endpoints : (the maximal edge of any cycle is vertical)



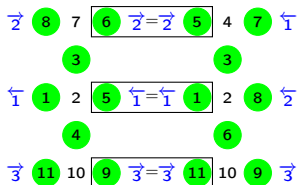
# Right compatibility relation



## Necessary and sufficient conditions

- ▶ the orders on common edges are the same
- ▶ horizontal edges are after at least one of their endpoints
- ▶ horizontal edges are before at least one of their endpoints : (the maximal edge of any cycle is vertical)

# Right compatibility relation



## Necessary and sufficient conditions

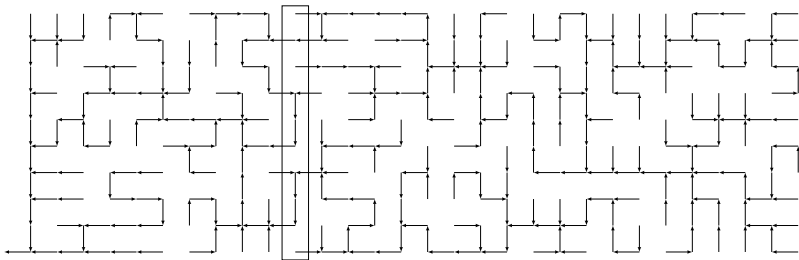
- ▶ the orders on common edges are the same
- ▶ horizontal edges are after at least one of their endpoints
- ▶ horizontal edges are before at least one of their endpoints : (the maximal edge of any cycle is vertical)

Reduction on the number of states : permutations and orientations on the right horizontal edges

- ▶  $\#states \leq 2^H \cdot H! (\leq (4H - 1)!)$
- ▶ the transitions and alphabet are still the induced orders

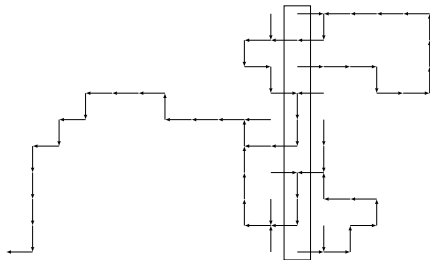
## Proposition

Any induced order appears as a letter on some recurrent configuration's standardized decomposition on  $[1, W] \times [1, H]$  for  $W = 2H + 1$ .



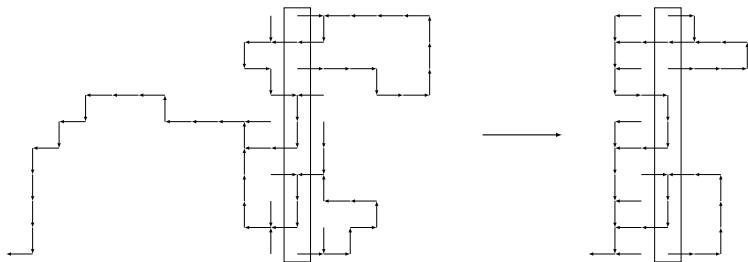
## Proposition

Any induced order appears as a letter on some recurrent configuration's standardized decomposition on  $[1, W] \times [1, H]$  for  $W = 2H + 1$ .



## Proposition

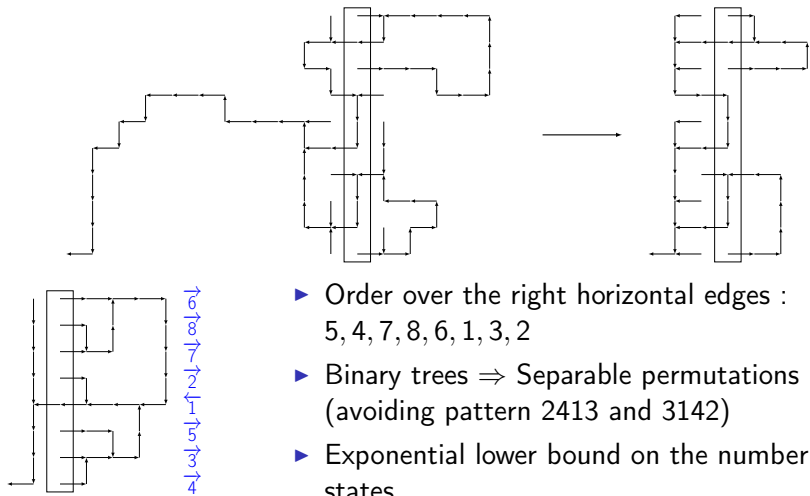
Any induced order appears as a letter on some recurrent configuration's standardized decomposition on  $[1, W] \times [1, H]$  for  $W = 2H + 1$ .





## Proposition

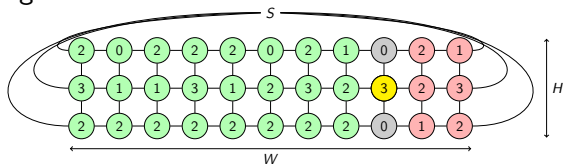
Any induced order appears as a letter on some recurrent configuration's standardized decomposition on  $[1, W] \times [1, H]$  for  $W = 2H + 1$ .



- ▶ Order over the right horizontal edges :  
5, 4, 7, 8, 6, 1, 3, 2
- ▶ Binary trees  $\Rightarrow$  Separable permutations  
(avoiding pattern 2413 and 3142)
- ▶ Exponential lower bound on the number of states

# Towards the infinite ladder (Járai, Lyons 07)

Approximation of  $\mathbf{Z} \times [1, H]$  by square grid graphs rooted to the sink on the left and right sides.



- ▶ vertices that topples from the left (green and yellow)
- ▶ vertices that topples from the right (red and yellow)
- ▶ other vertices (gray)

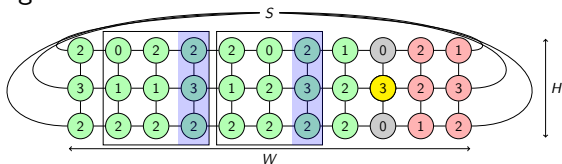
A *left burnable configuration* is a green/yellow configuration.

Decomposition in left-burnable blocks..

Automaton on a super set of column alphabet  $\Rightarrow$  Parry measure  $\Rightarrow$  Uniform distribution on left-burnable configurations of  $\mathbf{Z} \times [1, H]$

# Towards the infinite ladder (Járai, Lyons 07)

Approximation of  $\mathbf{Z} \times [1, H]$  by square grid graphs rooted to the sink on the left and right sides.



- ▶ vertices that topples from the left (green and yellow)
- ▶ vertices that topples from the right (red and yellow)
- ▶ other vertices (gray)

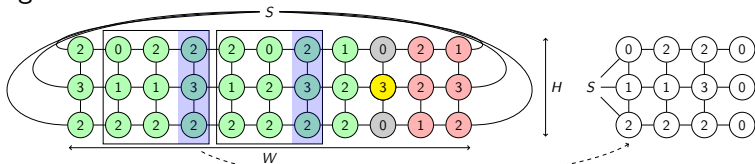
A *left burnable configuration* is a green/yellow configuration.

Decomposition in left-burnable blocks..

Automaton on a super set of column alphabet  $\Rightarrow$  Parry measure  $\Rightarrow$  Uniform distribution on left-burnable configurations of  $\mathbf{Z} \times [1, H]$

# Towards the infinite ladder (Járai, Lyons 07)

Approximation of  $\mathbf{Z} \times [1, H]$  by square grid graphs rooted to the sink on the left and right sides.



- ▶ vertices that topples from the left (green and yellow)
- ▶ vertices that topples from the right (red and yellow)
- ▶ other vertices (gray)

A *left burnable configuration* is a green/yellow configuration.

Decomposition in left-burnable blocks..

Automaton on a super set of column alphabet  $\Rightarrow$  Parry measure  $\Rightarrow$  Uniform distribution on left-burnable configurations of  $\mathbf{Z} \times [1, H]$

THANK YOU